

The Quantum Decoding Problem : Tight Achievability Bounds and Application to Regev's Reduction

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Linear codes

Consider \mathcal{C} an $[n, k]$ linear code on a finite field \mathbb{Z}_q i.e \mathbb{Z}_q -linear subspace of \mathbb{Z}_q^n of length n and dimension k . q is prime

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The decoding problem

Input $y = c + e \in \mathbb{Z}_q^n$ with $c \in \mathcal{C}$ and $e \in \mathbb{Z}_q^n$ an error chosen with distribution p .

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The Short Codeword Problem (SCP)

Input \mathcal{C} linear code $[n, k]$, a metric $|\cdot|$ on \mathbb{Z}_q^n and a bound ω

Output c a non-zero codeword verifying $|c| \leq \omega$

Why are we interested in these problems?

Uses in cryptography:

- Decoding Problem:
 - Encryption : Bike, McEliece, HQC, Kyber
 - Signature : SDitH
- Short Codeword Problem:
 - Signature : Wave , Dillithium

Idea : Study the difficulty to solve these problems with the quantum computer.

Quantum algorithm : quantum reduction SCP (multiple solutions) $<$ DP (one solution)

Solve SCP using :

- a classical algorithm solving the decoding problem
- the quantum fourier transform

General algorithm (Regev)

Classical Fourier transform : $\hat{f}(x) = \frac{1}{\sqrt{q^n}} \sum_{y \in \mathbb{Z}_q^n} e^{\frac{2i\pi \langle x, y \rangle}{q}} f(y)$

Quantum Fourier transform : $\text{QFT}|x\rangle = \frac{1}{\sqrt{q^n}} \sum_{y \in \mathbb{Z}_q^n} e^{\frac{2i\pi \langle x, y \rangle}{q}} |y\rangle$

We denote the **dual code** of \mathcal{C} as $\mathcal{C}^\perp = \{x \in \mathbb{Z}_q^n \mid x \cdot c^T = 0 \text{ for all } c \in \mathcal{C}\}$

Lemma

$$\text{QFT}\left(\sum_{c \in \mathcal{C}} \sum_e F(e) |c + e\rangle\right) = \sum_{c^\perp \in \mathcal{C}^\perp} \hat{F}(c^\perp) |c^\perp\rangle$$

General algorithm (Regev)

1. *Construct* $\frac{1}{\sqrt{q^k}} \sum_{c \in C} |c\rangle \otimes \sum_e F(e)|e\rangle$
2. *Intricate* $\sum_{c \in C} \sum_e F(e)|c\rangle \otimes |c + e\rangle$
3. *Decode* Quantumly solve the decoding problem : from $c + e$, find c
 $\frac{1}{Z} \sum_{c \in C} \sum_e F(e)|0\rangle|c + e\rangle.$
4. *Quantum Fourier Transform* $\frac{1}{Z} \sum_{c^\perp \in C^\perp} \hat{F}(c^\perp)|0\rangle|c^\perp\rangle$
5. *Measure*

Regev's algorithm modified (Chen, Liu, Zhandry)

1. *Construct* $\frac{1}{\sqrt{q^k}} \sum_{c \in C} |c\rangle \otimes \sum_e F(e)|e\rangle$
2. *Intricate* $\sum_{c \in C} \sum_e F(e)|c\rangle \otimes |c + e\rangle = \sum_{c \in C} |c\rangle |\psi_c\rangle$ with $|\psi_c\rangle = \sum_{e \in \mathbb{Z}_q^n} F(e)|c + e\rangle$
3. *Decode* Quantumly solve the decoding problem : from $|\psi_c\rangle$, find c
 $\frac{1}{Z} \sum_{c \in C} |0\rangle |\psi_c\rangle.$
4. *Quantum Fourier Transform* $\frac{1}{Z} \sum_{c^\perp \in C^\perp} \hat{F}(c^\perp) |0\rangle |c^\perp\rangle$
5. *Measure*

The quantum decoding problem

The quantum decoding problem (S-LWE)

Input $|\psi_c\rangle = \sum_{e \in \mathbb{Z}_q^n} F(e)|c + e\rangle$ with $e \in \mathbb{Z}_q^n$ an error and $F(e) = \sqrt{p(e)}$

Output c

The quantum decoding problem is not harder than the classical decoding problem

When we measure $|\psi_c\rangle$, we get $c + e$ noisy codeword with probability $|F(e)|^2$

[Chen, Liu, Zhandry 22]

Introduce **S-LWE** (quantum decoding problem): Quantum algorithms in **polynomial time** to solve **S-LWE** and a specific instance of **SIS[∞]** (SCP with infinite norm)

- particular zone of parameters : no classical algorithm doing this in **polynomial time**
- **reduction** from the quantum algorithm solving **QDP** to the quantum algorithm solving **SCP**
- not in a constant rate : **tends to 0**

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The quantum decoding problem [Chailloux, Tillich 23]

- natural zone of parameters: **constant rate**
- Bernoulli distribution
- solve the quantum decoding problem in polynomial time
- quantum algorithm in polynomial time to solve **SCP** equivalent to classical Prange algorithm
- go beyond with a **non polynomial time** algorithm to find minimal weight codeword in the dual code (PGM)

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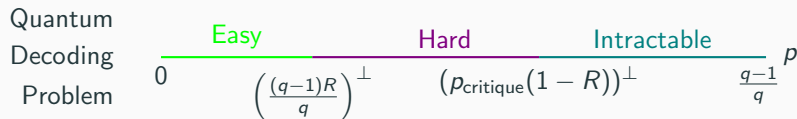
No exponential quantum speedup for **SIS[∞]** anymore [Kothary, O'Donnell, Wu 25]

- **Classical algorithm** that solves **SIS[∞]** in polynomial time
- No exponential quantum speedup anymore

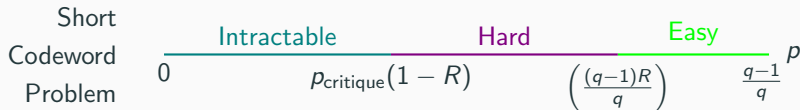
Why is it interesting?



$$\Pr(e_i = 0) = 1 - p$$



$$\Pr(e_i \neq 0) = \frac{p}{q-1}$$



$$p^{\perp} = \frac{\sqrt{(1-p)(q-1)} - \sqrt{p}}{q}$$

The Pretty Good Measurement

Problem : From an ensemble $\{|\psi_i\rangle\}_{i\in[n]}$ of quantum states , we want to recover i from $|\psi_i\rangle$ when i is chosen at random.

Definition

The PGM associated to the ensemble $\{|\psi_i\rangle\}_{i\in[n]}$ of quantum states is the POVM $\{M_i\}_{i\in[n]}$ with

$$M_i = \rho^{-\frac{1}{2}} |\psi_i\rangle \langle \psi_i| \rho^{-\frac{1}{2}} \text{ given } \rho = \sum_{i\in[n]} |\psi_i\rangle \langle \psi_i|$$

Proposition

The PGM is optimal

Achievability result for random linear codes

We analyse the PGM in a general case of noise distribution $f = g^{\otimes n}$

Let $R = \frac{k}{n}$ be the rate of the code \mathcal{C} and $H_q(|\hat{g}|^2) = - \sum_{e \in \mathbb{Z}_q} |\hat{g}(e)|^2 \log(|\hat{g}(e)|^2)$ be the entropy.

Theorem : Achievability result for random linear codes [[Blanvillain, Chailloux, Tillich](#)]

If there exists $\delta > 0$ such that $R < H_q(|\hat{g}|^2) + \delta$, then

$$P_{\text{PGM}} = 1 - o(1)$$

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Theorem : Tractability for Hamming distribution [[Chailloux, Tillich](#)]

Let $g(e)$ the error function for the Bernoulli distribution such that

$$g(0) = \sqrt{1-p} \text{ and for } e \in \mathbb{Z}_q \setminus \{0\}, g(e) = \sqrt{\frac{p}{p-1}}.$$

If $p < (\delta_{\min}(1-R))^\perp$, then $P_{\text{PGM}} = 1 - o(1)$

Non-achievability result in the general case

Theorem : Achievability result for random linear codes [[Blanvillain, Chailloux, Tillich](#)]

If there exists $\delta > 0$ such that $R < H_q(|\hat{g}|^2) - \delta$, then

$$P_{\text{PGM}} = 1 - o(1)$$

Theorem : Non-achievability result in the general case [[Blanvillain, Chailloux, Tillich](#)]

If there exists $\delta > 0$ such that $R \geq H_q(|\hat{g}|^2) + \delta$, then for any quantum algorithm and any code with q^{Rn} codewords,

$$P_{\text{PGM}} = o(1)$$

The rank metric case

Rank metric

When $n = a \cdot b$, $a \geq b$

Arrange the entries of a vector $\mathbf{x} \in \mathbb{Z}_q^n$ in a matrix $\mathbf{X} = \text{Mat}(\mathbf{x}) \in \mathbb{Z}_q^{a \times b}$

$$|\mathbf{x}|_{\text{rk}} \triangleq \text{rank}(\mathbf{X})$$

$$f_t^{a,b}(\mathbf{e}) = \begin{cases} \frac{\begin{bmatrix} b - |\mathbf{e}|_{\text{rk}} \\ t - |\mathbf{e}|_{\text{rk}} \end{bmatrix}_q}{\sqrt{q^{at}Z}} & \text{if } |\mathbf{e}|_{\text{rk}} \leq t \\ 0 & \text{else} \end{cases} \quad \text{with} \quad \begin{bmatrix} b \\ t \end{bmatrix}_q = \begin{cases} \prod_{i=0}^{t-1} \frac{q^b - q^i}{q^t - q^i} & \text{if } t \leq b \\ 0 & \text{else} \end{cases}$$

$$\hat{f}_t = f_{b-t}. \tag{1}$$

Application in rank metric : case of interest in code-based cryptography : alternative to the usual Hamming metric

The distribution does **not** correspond to a product distribution

The probability distribution is a decreasing function of the rank weight. We get at the limit where the PGM works elements in the dual code of minimum rank weight.

The Quantum Decoding Problem : Takeaway

👉 **Tight achievability results** : it is possible from an information theoretic perspective to solve the Quantum Decoding Problem for random linear codes when $R < H_q(|\hat{g}|^2) - \delta$ for $\delta > 0$

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Thank you for your attention !