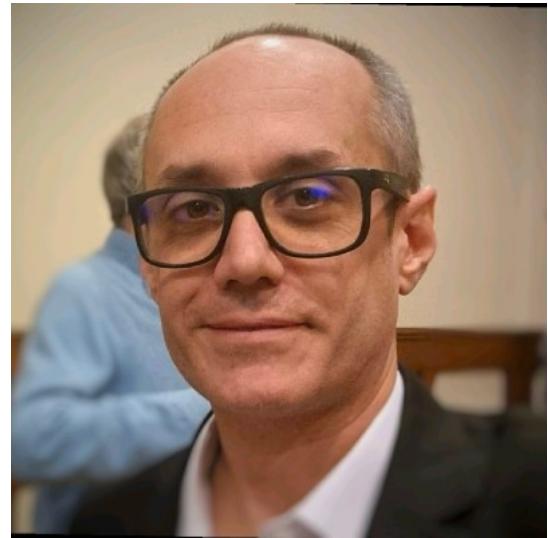


Verifiable blind observable estimation

Bo Yang, Elham Kashefi, Harold Ollivier

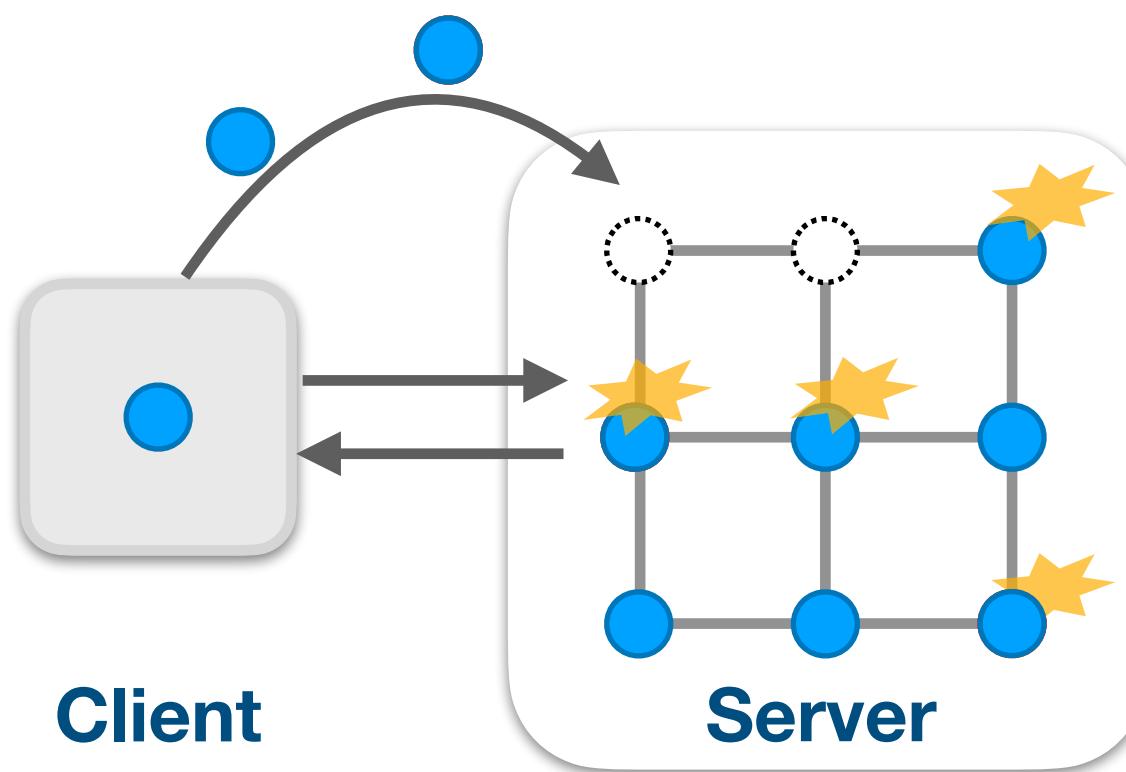
<https://www.arxiv.org/abs/2510.08548>



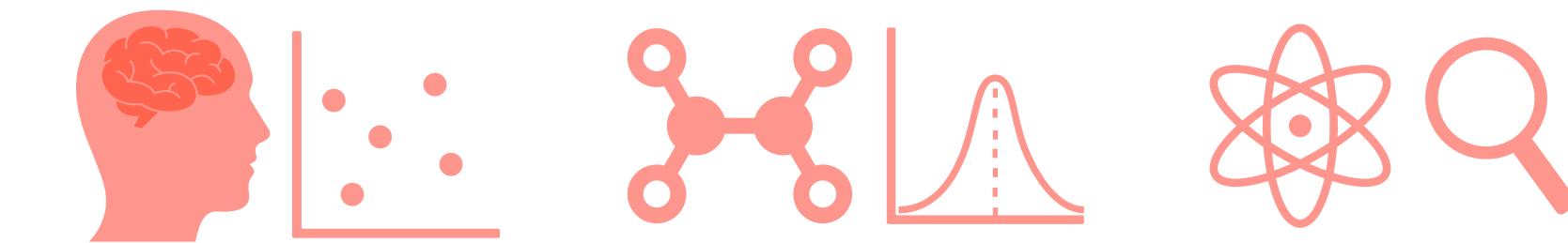
JIQ2026, Bordeaux, 16.01.2026

Our contribution

- The first **ideal resource** in abstract cryptography
- Overhead-free **secure verification protocol** for **observable estimation tasks**



Secure certification



QML / simulation / metrology

key quantum applications

Verification Protocols

Observable Estimation

Towards **secure verifiable** near-term quantum advantage

Towards practical quantum utility

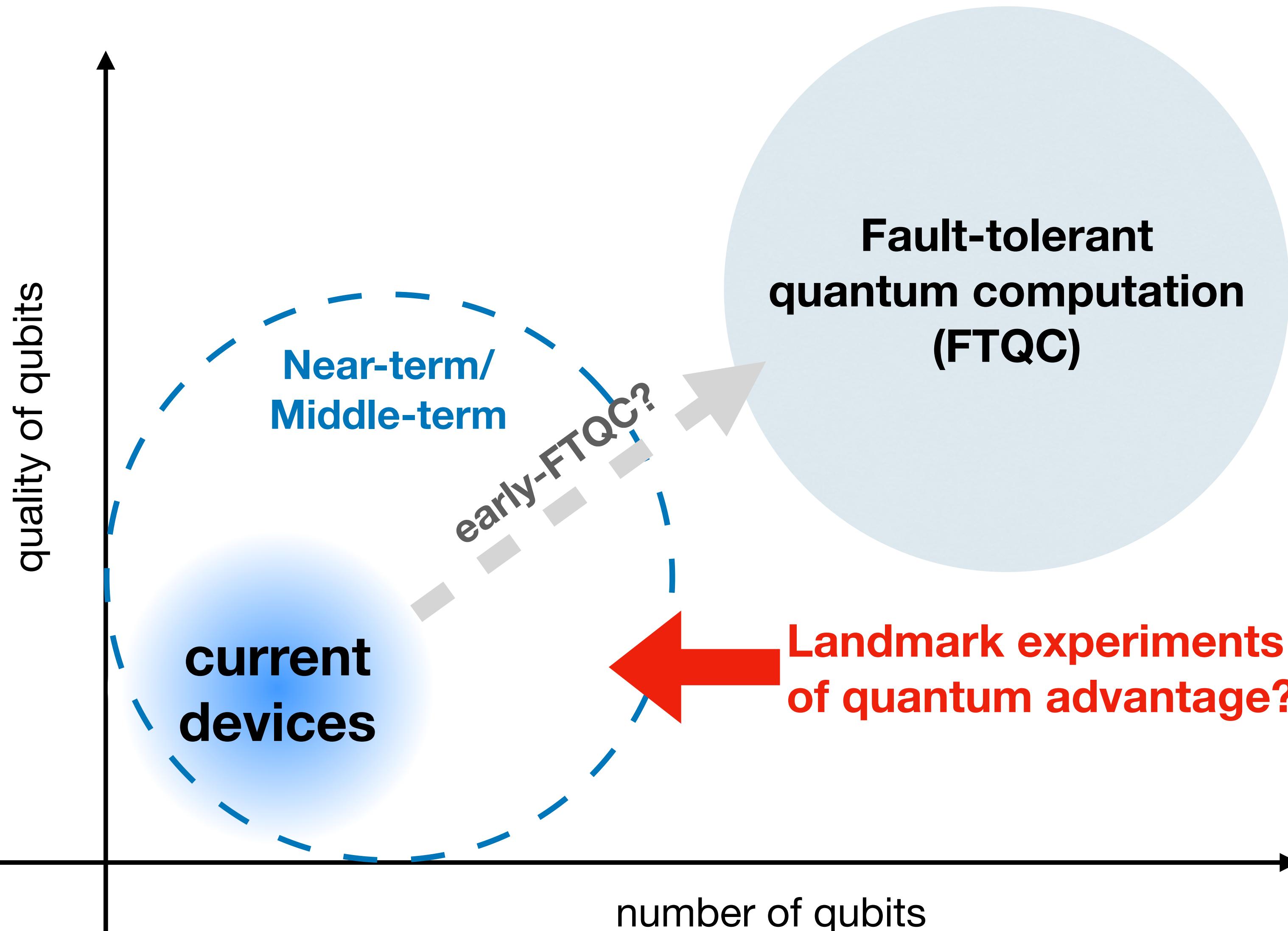
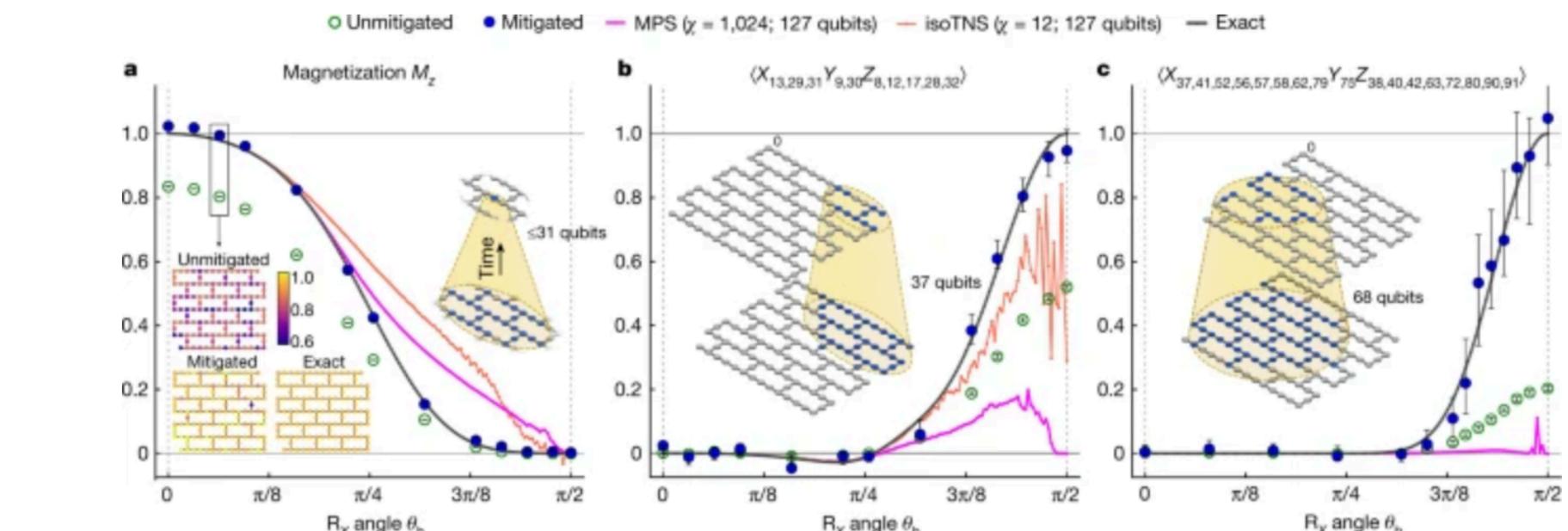
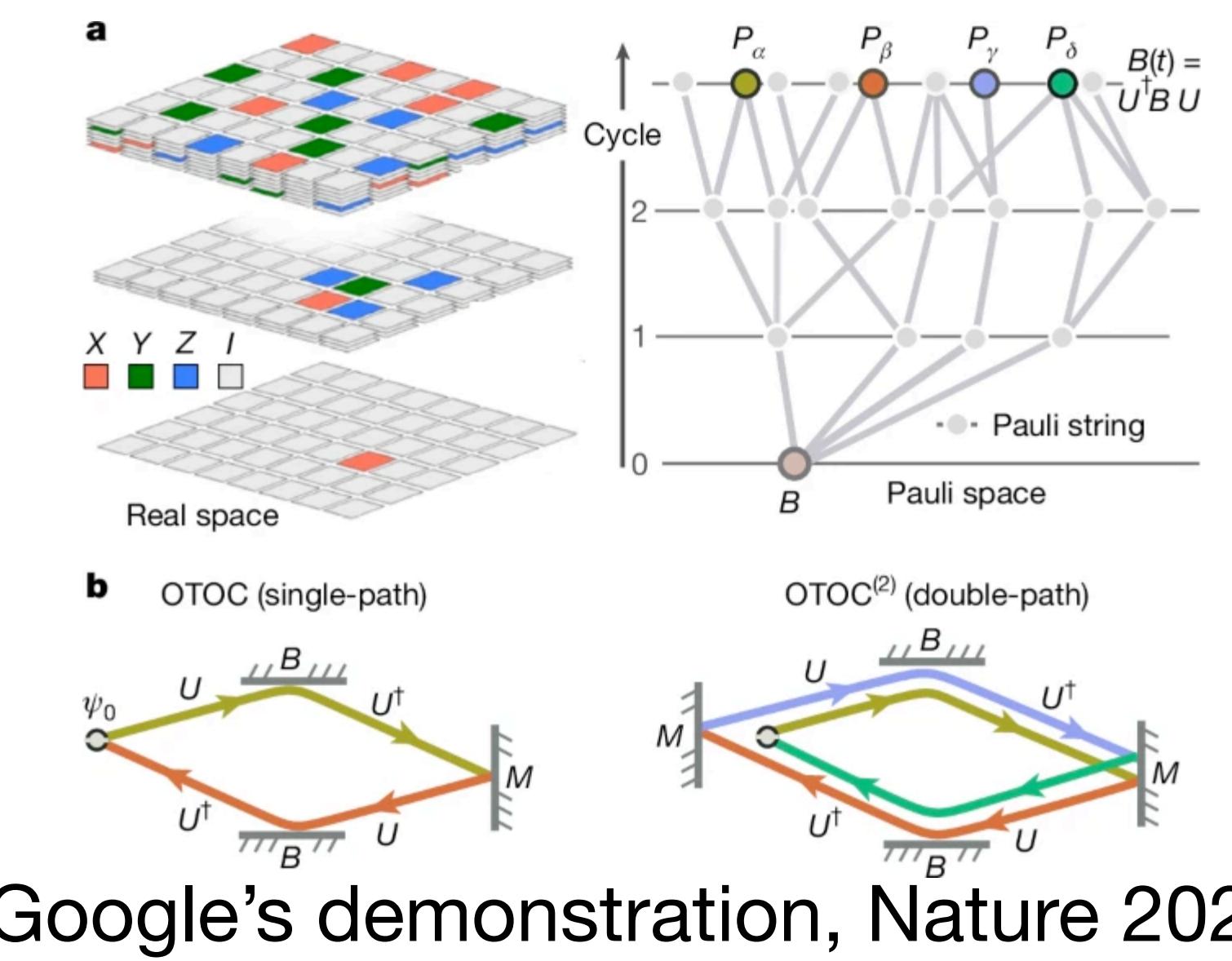


Fig. 3: Classically verifiable expectation values from 127-qubit, depth-15 Clifford and non-Clifford circuits.



[IBM's demonstration, Nature 2024]

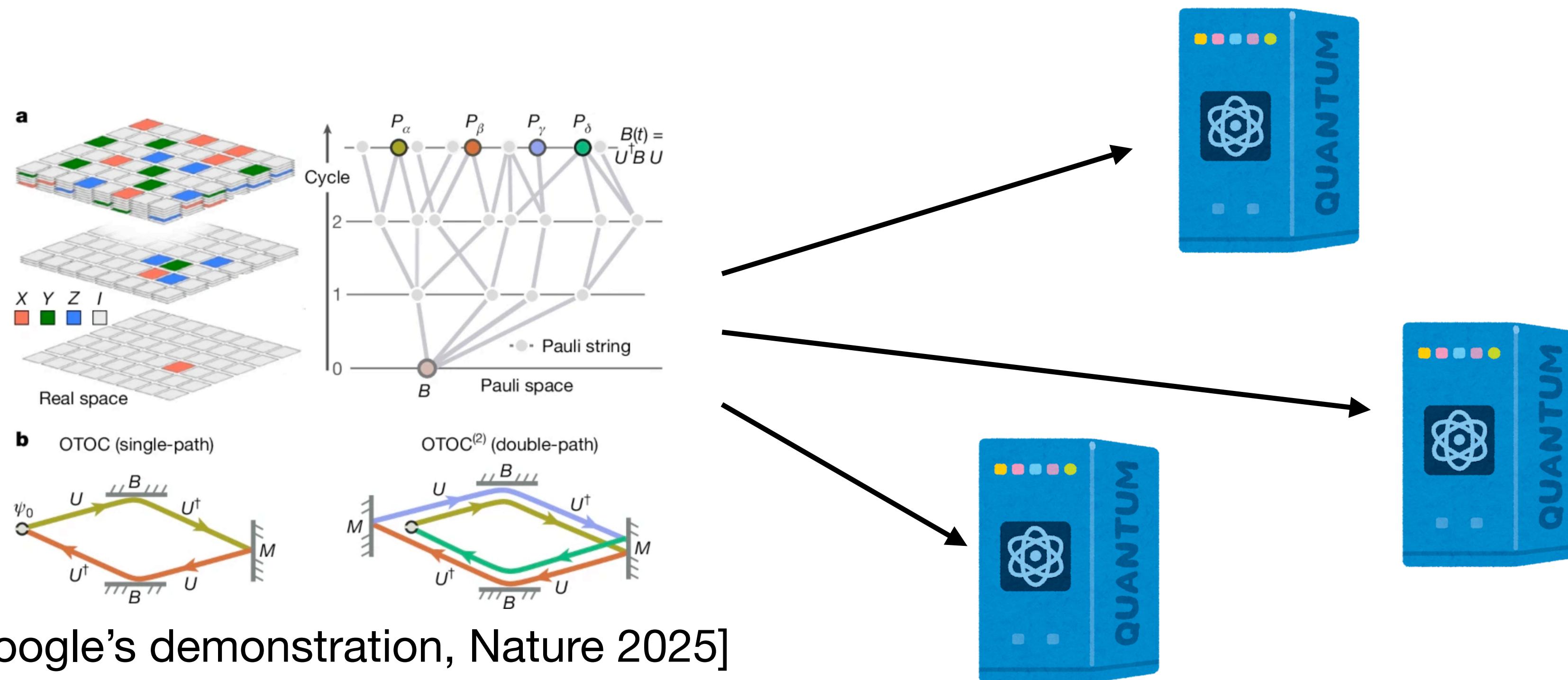


[Google's demonstration, Nature 2025]

Verifiability of quantum advantage

Definition of quantum advantage [Lanes+25]

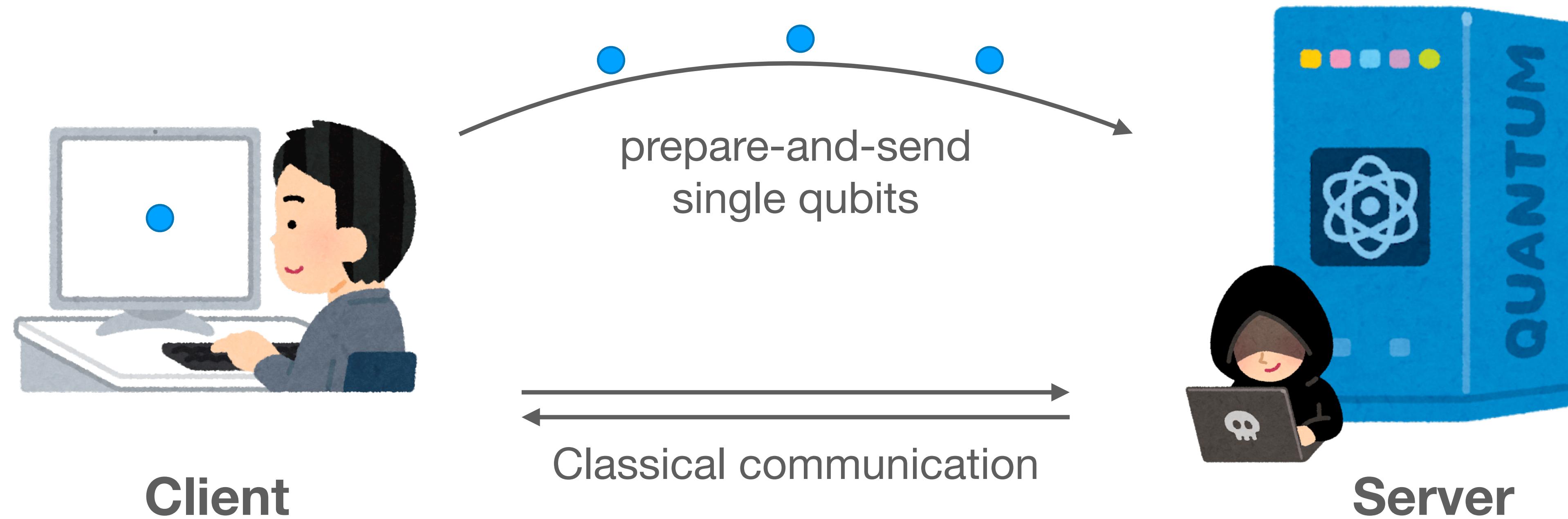
1. demonstrably superior to classical computational resources
2. **outputs can be rigorously validated**



Q. How to verify the outcome of classically intractable problem?

Secure delegated quantum computation

→ to untrusted and unreliable server

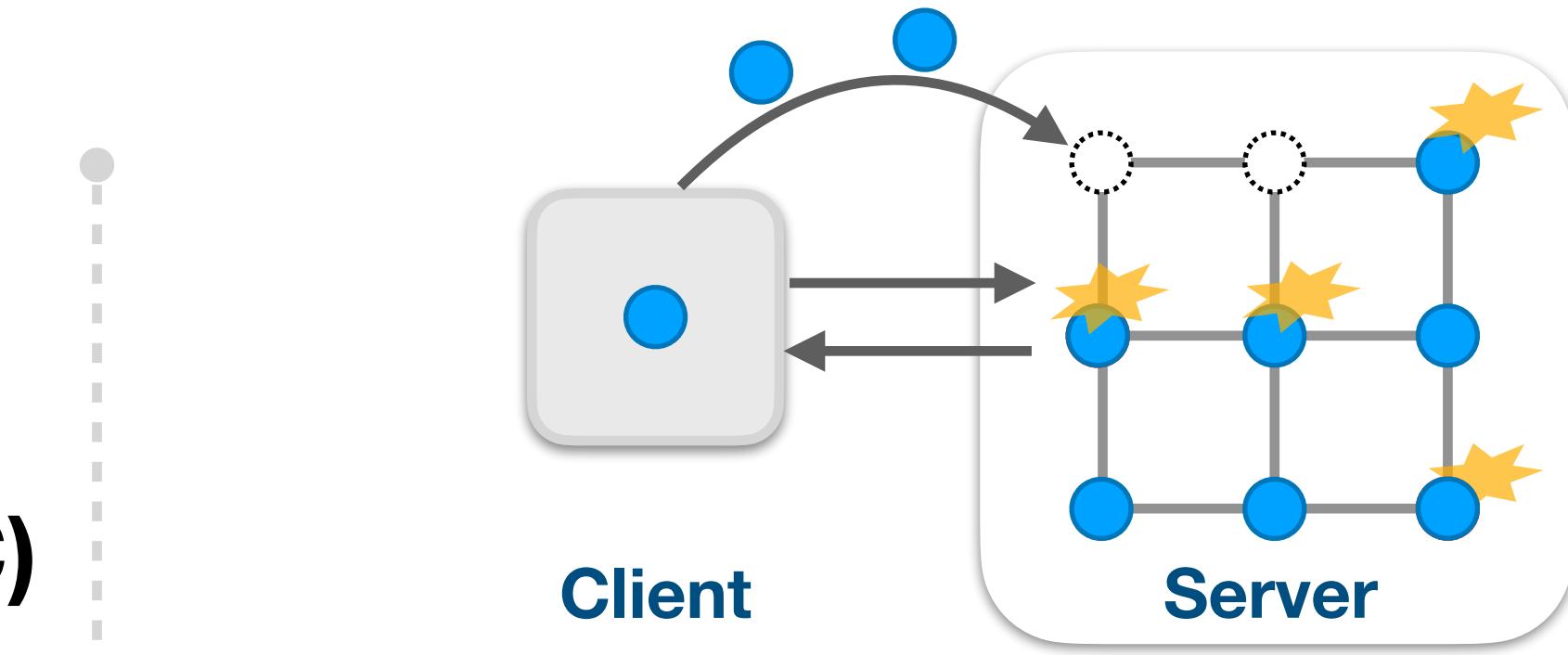


How to ensure the **integrity** and **confidentiality** of
delegated quantum computation?

**Universal Blind
Quantum Computation
(UBQC) [BFK09]**

- Perfect blindness

**Perfect security in
abstract cryptography (AC)**



The Client can prepare-and-send
single qubits

Blindness:

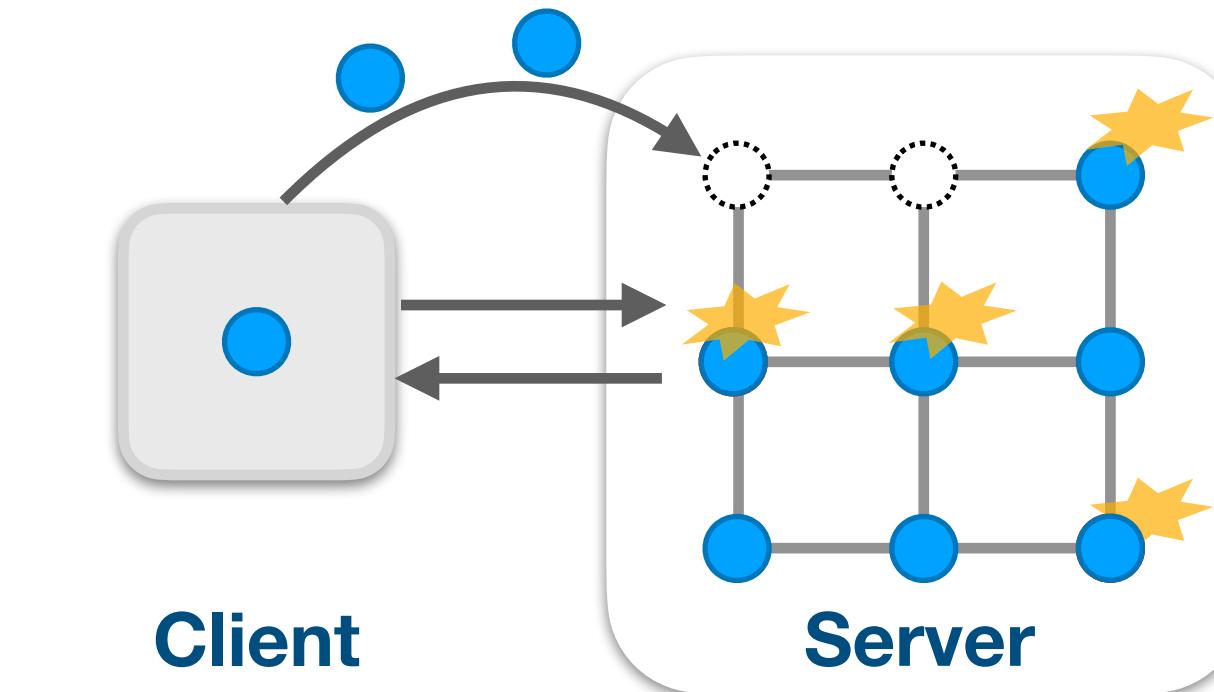
secret parameters in \bullet , $| +_\theta \rangle$,

Only Client can decode them
within the MBQC process

Universal Blind Quantum Computation (UBQC) [BFK09]

- Perfect blindness

Perfect AC security

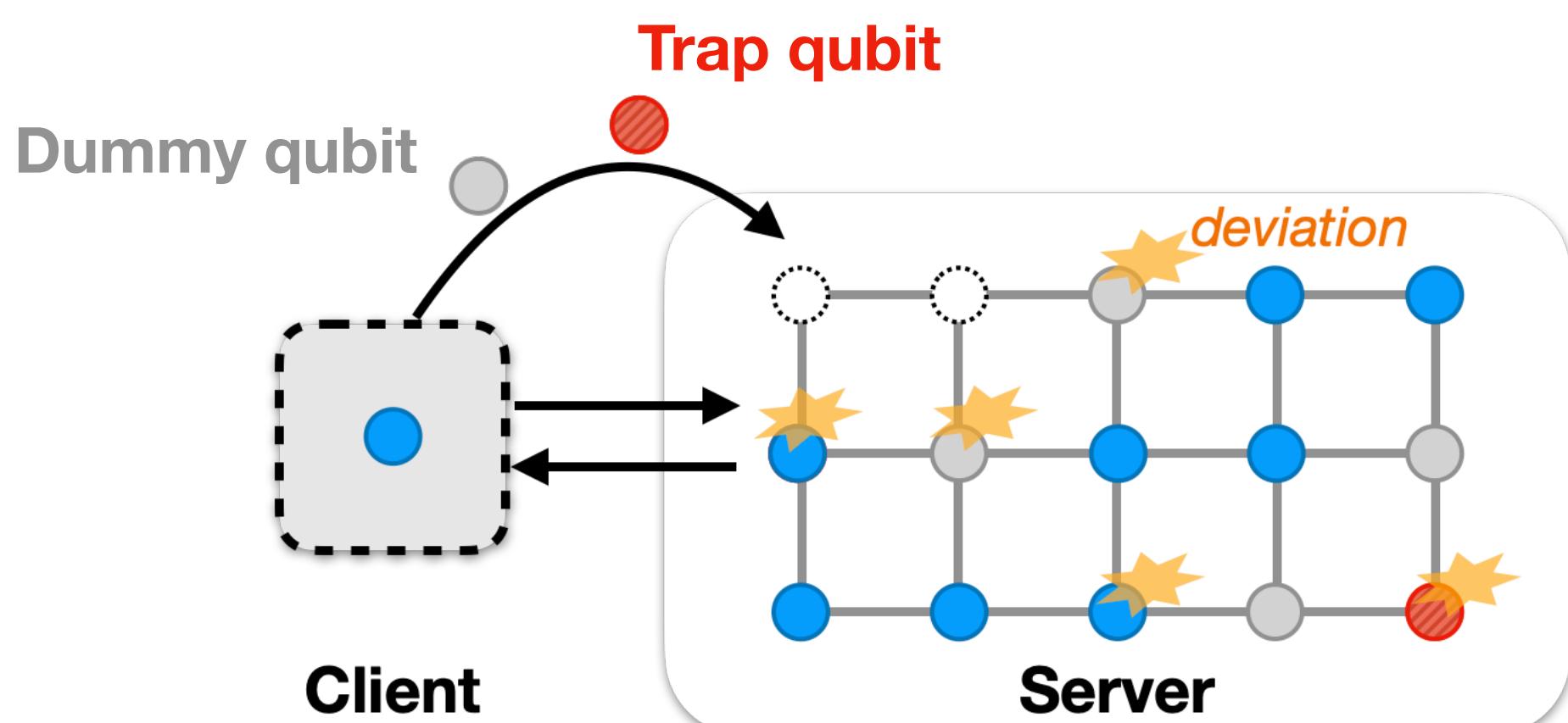


Verifiable Blind Quantum Computation (VBQC) [FK12]

- Perfect blindness
- Exponential verifiability

Exponential AC security

- Embed **traps** into MBQC
- **Require fault-tolerance for security amplification**
- Quantum output



dummy $|d\rangle$ $d \in \{0,1\}$

trap $|+_\theta\rangle = \frac{1}{\sqrt{2}} (|0\rangle + e^{i\theta}|1\rangle)$

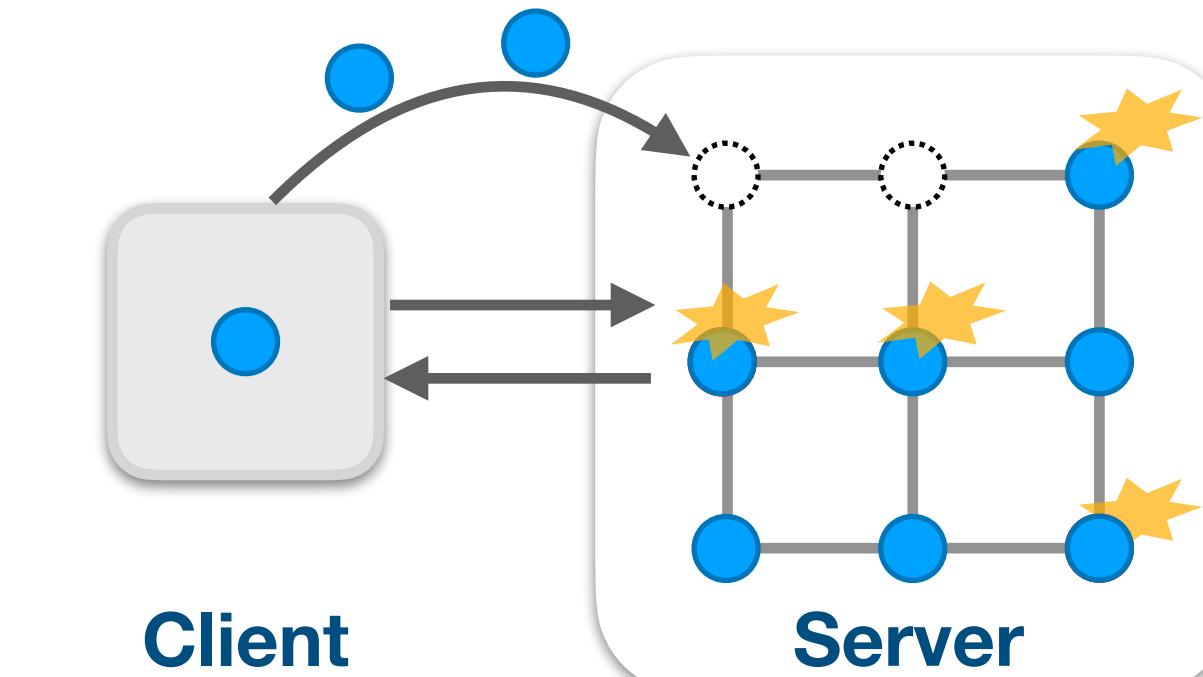
$$\text{CZ } |d\rangle |+_\theta\rangle = |d\rangle Z^d |+_\theta\rangle$$

Measurement basis $\{|+_\theta\rangle, |-_\theta\rangle\}$

Universal Blind Quantum Computation (UBQC) [BFK09]

- Perfect blindness

Perfect AC security



Verifiable Blind Quantum Computation (VBQC) [FK12]

- Perfect blindness
- Exponential verifiability

Exponential AC security

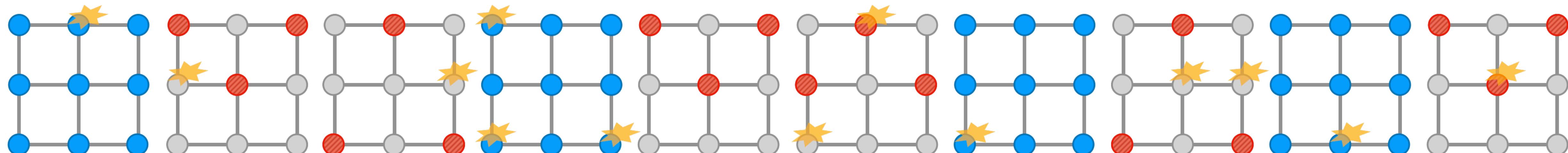
- Embed traps into MBQC
- **Require fault-tolerance for security amplification**
- Quantum output

Robust VBQC (RVBQC) [LMKO21]

- Perfect blindness
- Exponential verifiability

Exponential AC security

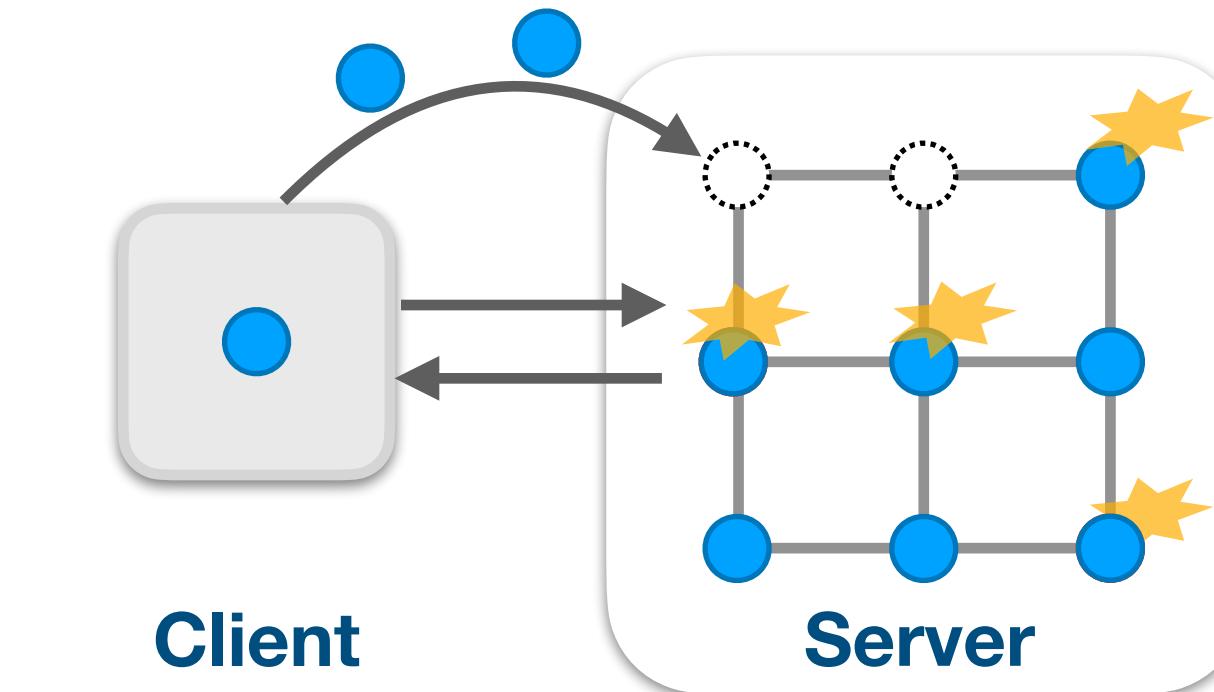
- Separated test/target runs
- **BQP decision problems**



Universal Blind Quantum Computation (UBQC) [BFK09]

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Verifiable Blind Quantum Computation (VBQC) [FK12]

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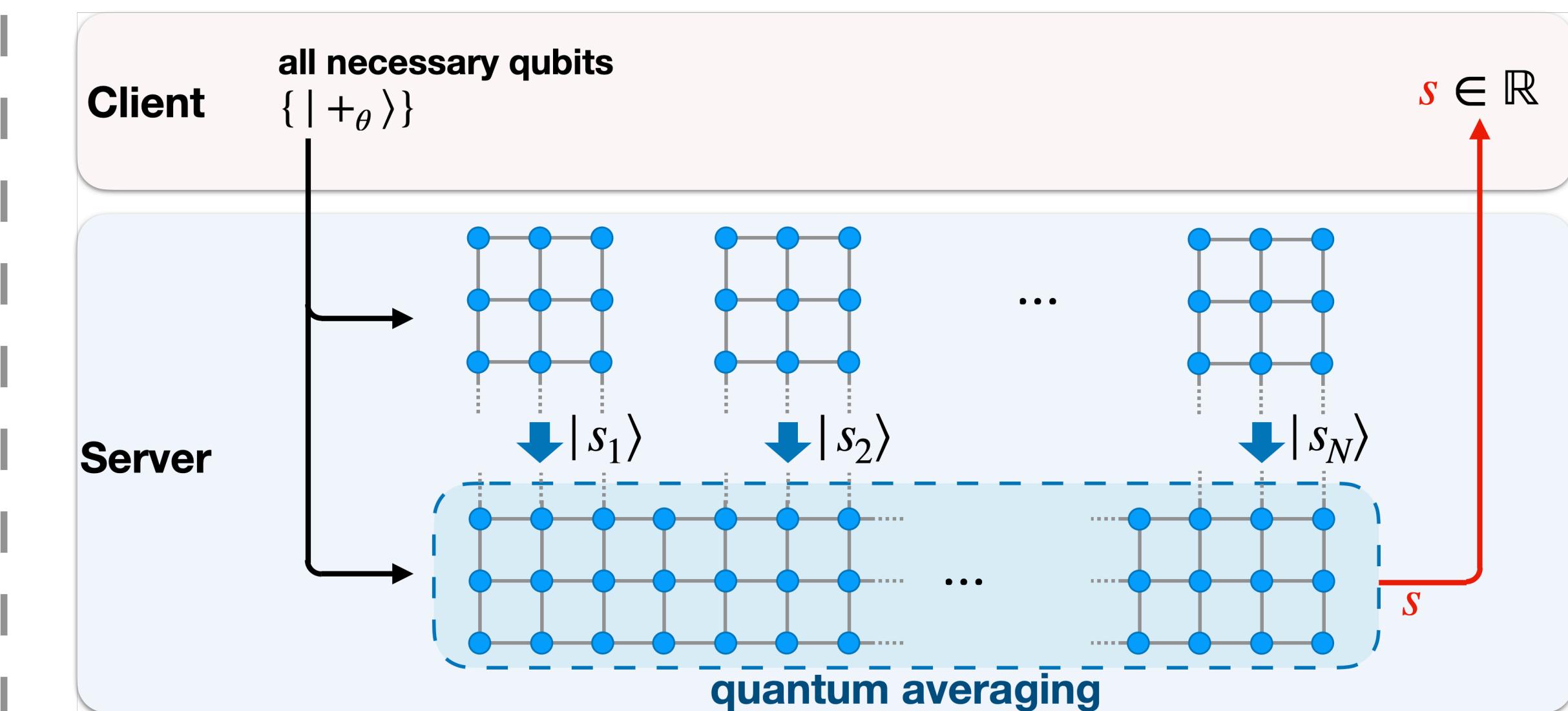
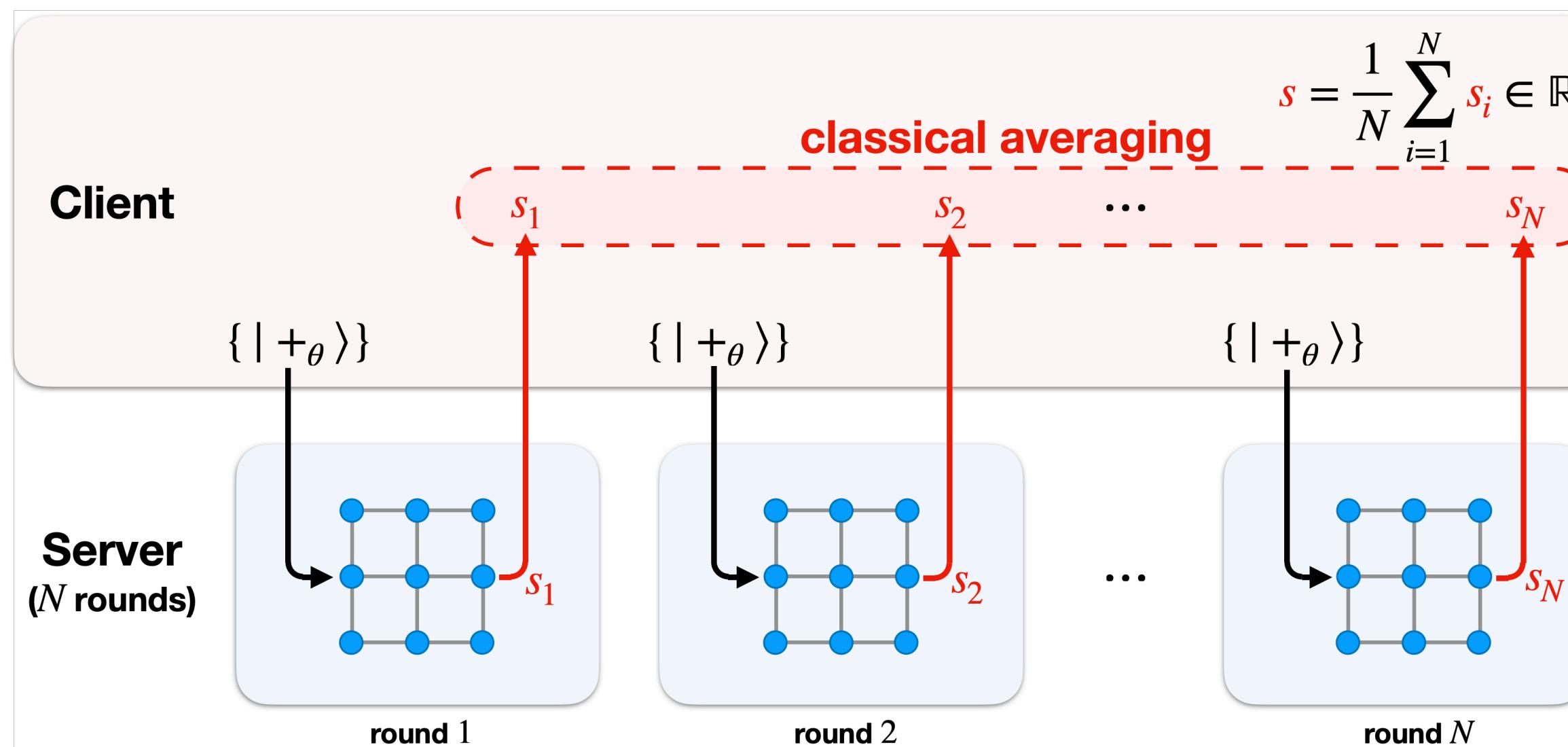
- Separated test/target runs
- **BQP decision problems**

However...

Verifying **continuous-valued output** with RVBQC requires **additional circuit overhead**.

What is the issue?

→ Either inverse polynomial security or heavy space overhead



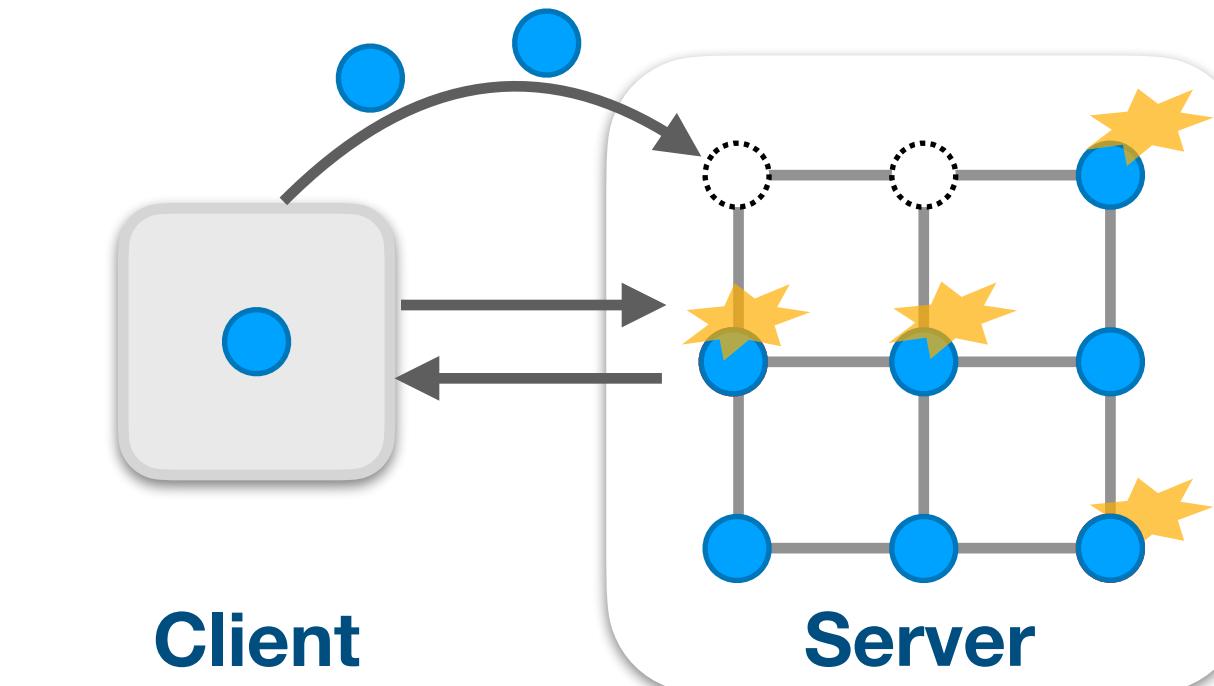
Only $\frac{N}{\text{poly}(n)}$ security
with conventional security notion

Beyond near-term quantum overhead
to achieve $O(e^{-N})$ security

Universal Blind Quantum Computation (UBQC) [BFK09]

- Perfect blindness

Perfect AC security



Verifiable Blind Quantum Computation (VBQC) [FK12]

- Perfect blindness
- Exponential verifiability

Exponential AC security

- Embed **traps** into MBQC
- **Require QEC for security**
- Quantum output

Robust VBQC (RVBQC) [LMKO21]

- Perfect blindness
- Exponential verifiability

Exponential AC security

- Separated test/target runs
- **BQP decision problems**

Verifiable Blind Observable Estimation (VBOE) [YKO25]

- Perfect **blindness**
- Exponential **verifiability**

Exponential AC security

- Separated test/target runs
- **Observable estimation**

Task: observable estimation

Given: Observable O , number of samples N , and allowed estimation bias ϵ .

Target: To estimate $\text{Tr} [\rho O]$ for a reference state ρ .

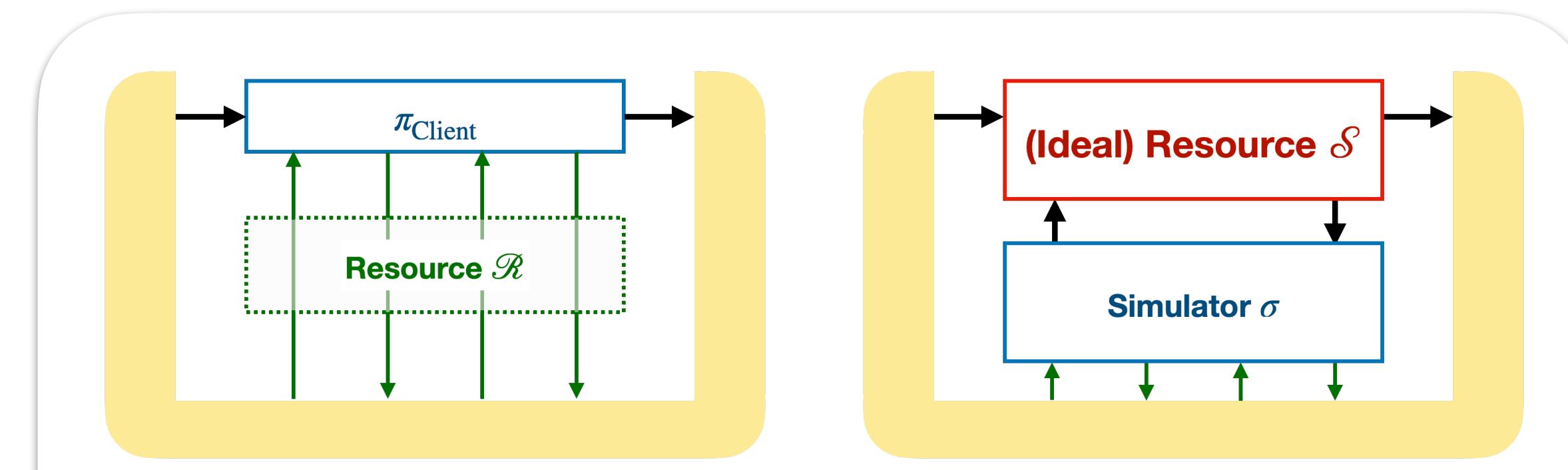
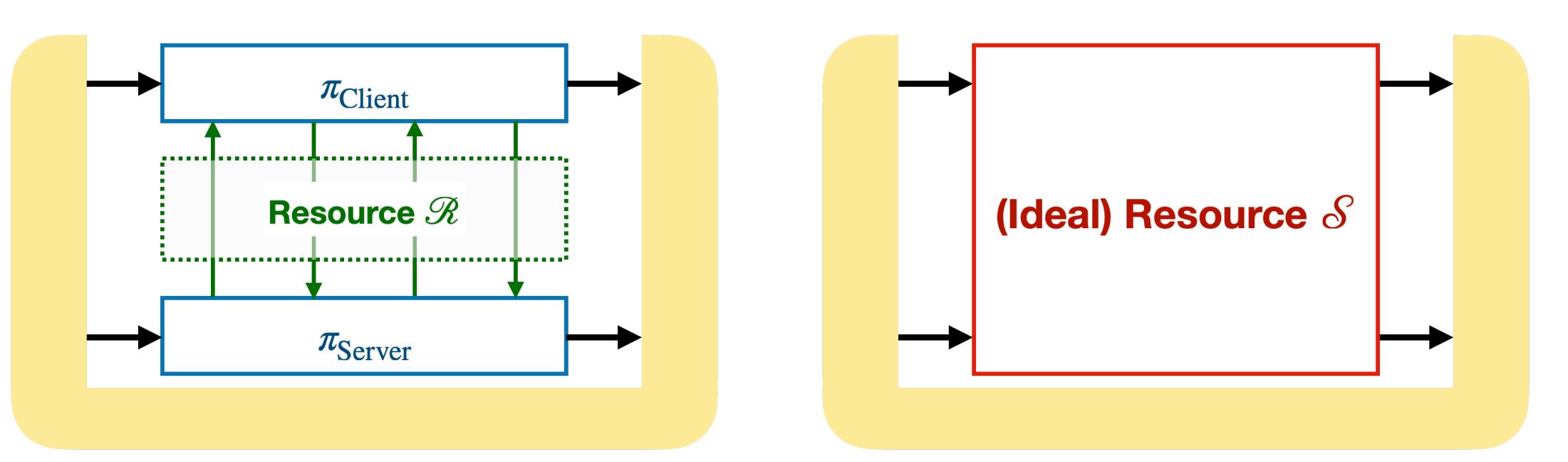
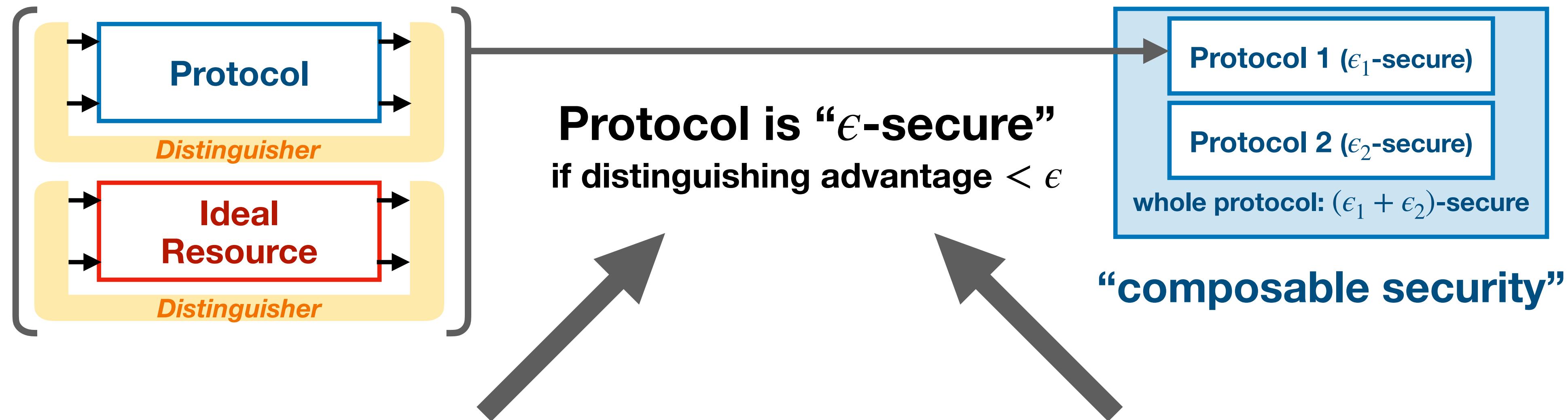
Output: Empirical mean $\hat{o} = \frac{1}{N} \sum_{i=1}^N \hat{y}_i$, where $\hat{y}_i \in \{-1,1\}$.

Performance: Estimation confidence $\Pr \left[\left| \hat{o} - \text{Tr} [\rho O] \right| \geq \epsilon \right] \leq \delta$.

→ “A protocol (ϵ, δ) -estimates $\text{Tr} [\rho O]$ ”.

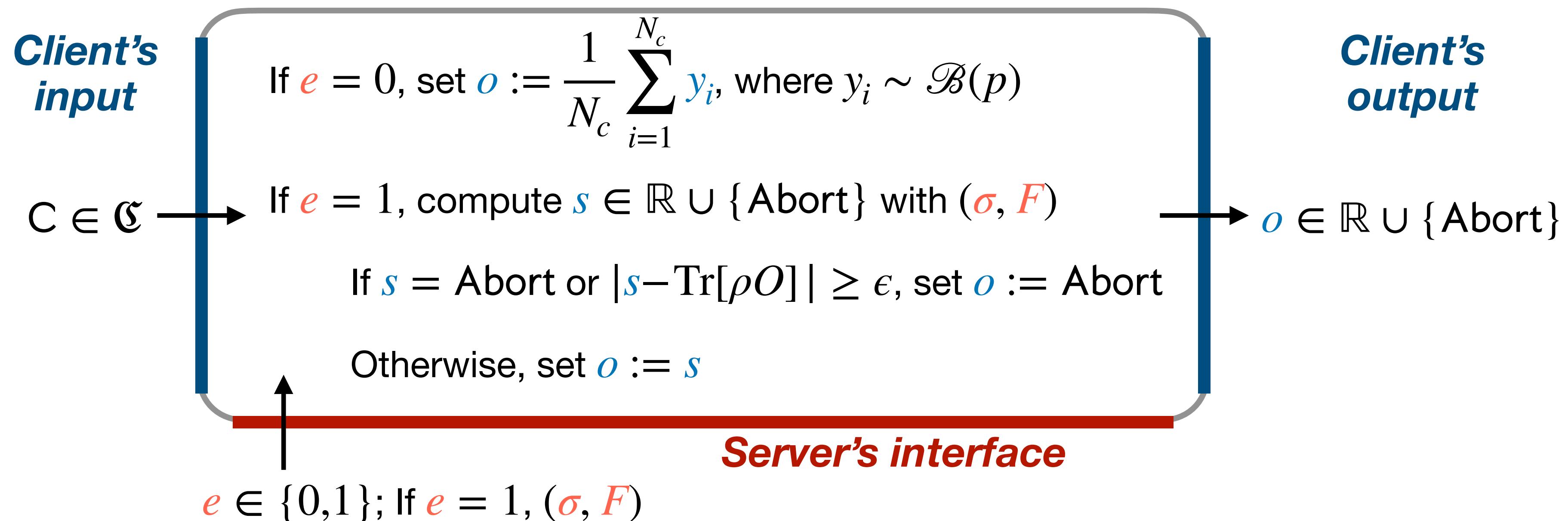
Without loss of generality,
 $O := |+\rangle\langle +| - |-\rangle\langle -|$

Abstract cryptography



Proposed ideal resource

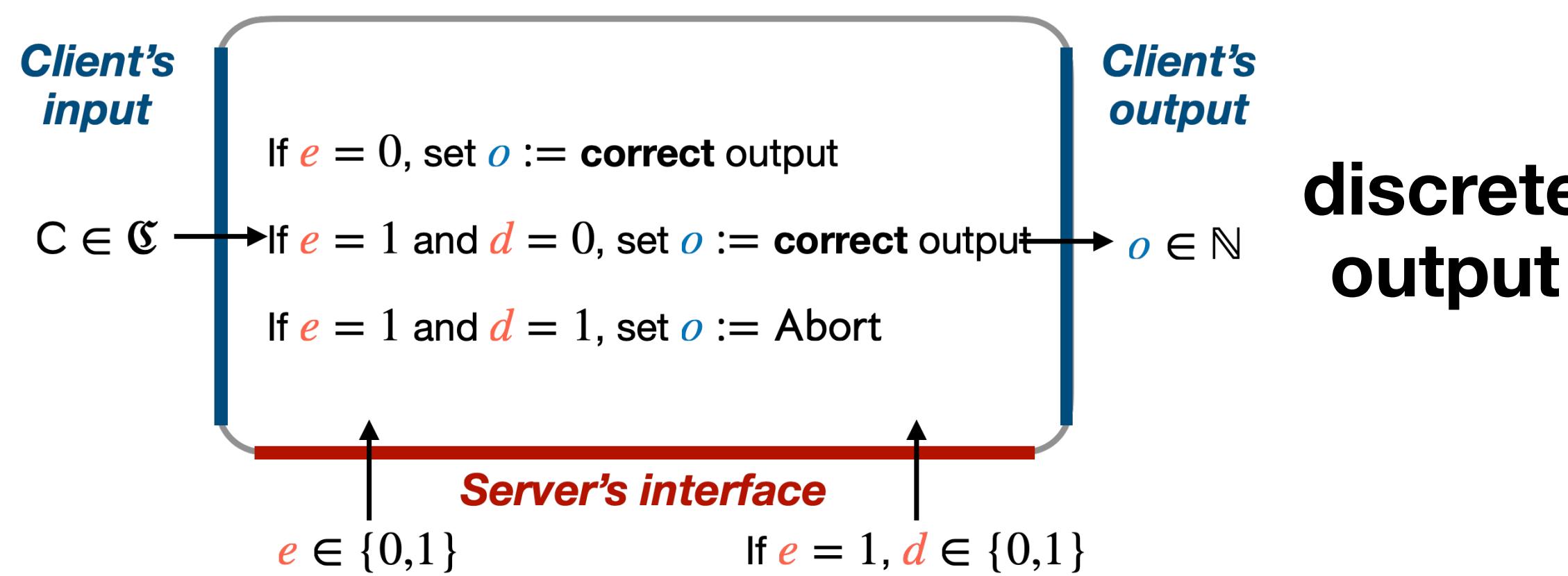
Secure Delegated Observable Estimation (SDOE)



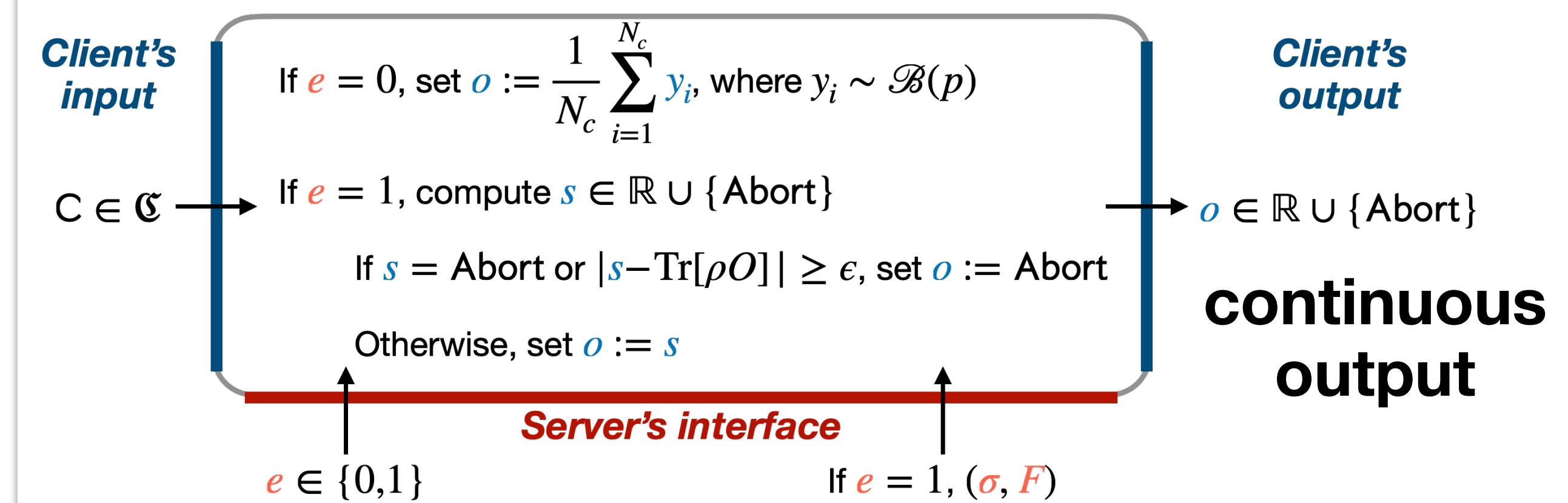
→ When accepted, the SDOE resource $(\epsilon, 0)$ -estimates $\text{Tr} [\rho O]$.

New conceptual tool in abstract cryptography

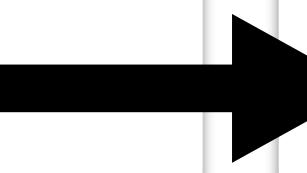
SDQC Resource [KKLO22], [DFPR14]



SDOE Resource [This work]



Strictly correct discrete output

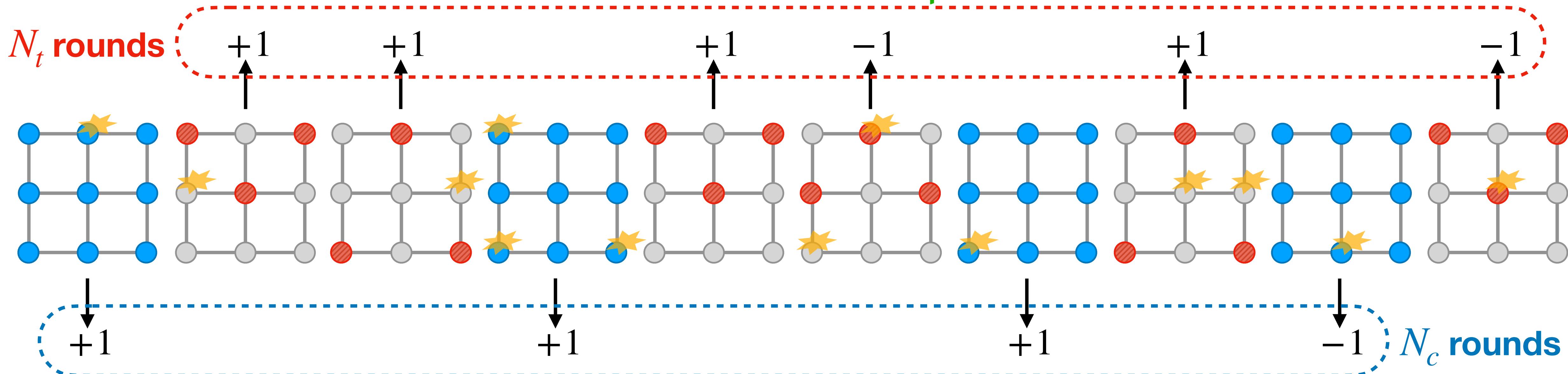


Allowing **biased** output (up to an allowed extent)

Proposed secure protocol

Verifiable Blind Observable Estimation (VBOE)

Test rounds: if satisfied, **Accept**, otherwise, return **Abort**.

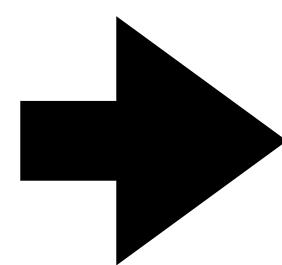


Computation rounds: If **Accept**, return $\tilde{o} = \frac{1}{N_c} \sum_{i=1}^{N_c} \tilde{y}_i$ to estimate $\text{Tr}[\rho O]$.

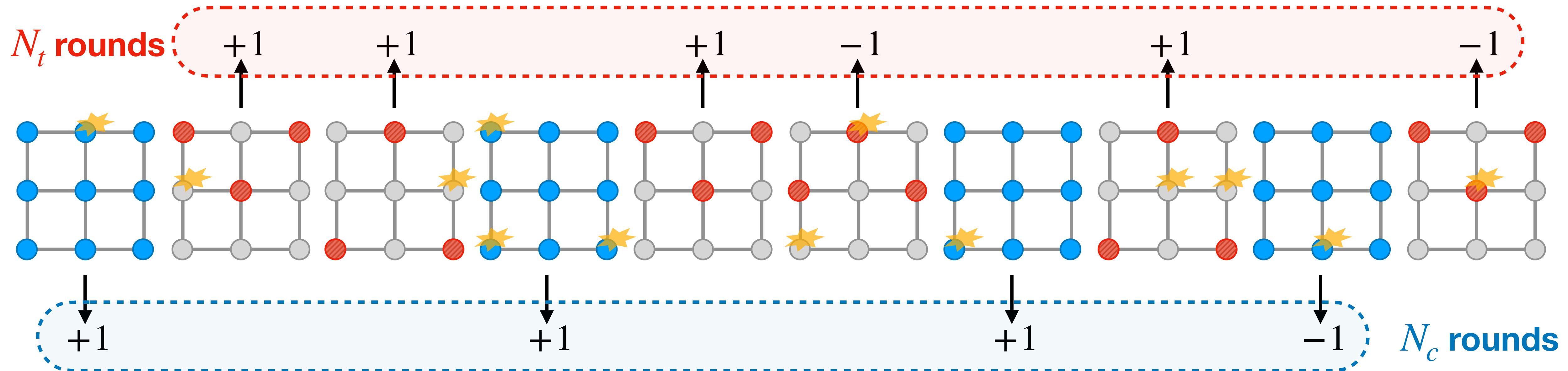
Main theorem

Exponential composable security of VBOE

VBOE (ϵ, δ) -estimates $\text{Tr} [\rho O]$ with exponentially small δ to N_c and N_t .

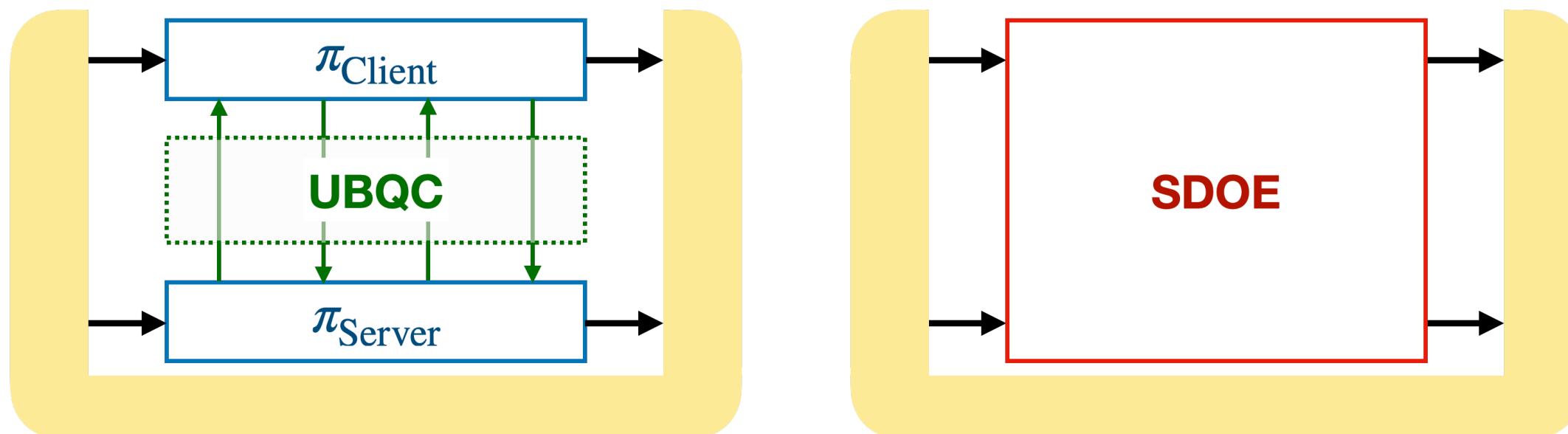


Exponentially small distinguishing advantage from SDOE,
i.e. $\delta = O(e^{-N_c} + e^{-N_t})$.



Proof sketch

Correctness (vs **honest, noise-free** Server)

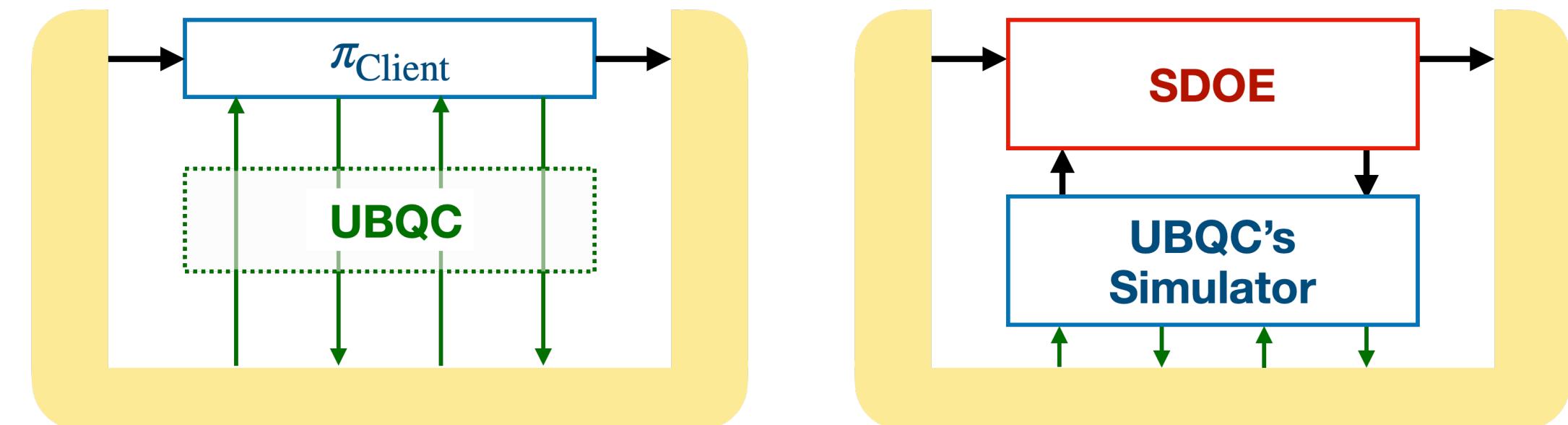


- VBOE always accepts
- SDOE further checks $|\hat{o} - \text{Tr} [\rho O]| \geq \epsilon$

↓ Difference in Accept/Abort probability

$$\Pr \left[|\hat{o} - \text{Tr} [\rho O]| \geq \epsilon \right] \leq 2 \exp \left(-\frac{\epsilon^2}{2} N_c \right)$$

Security (vs **malicious** Server)



- Using composability of UBQC

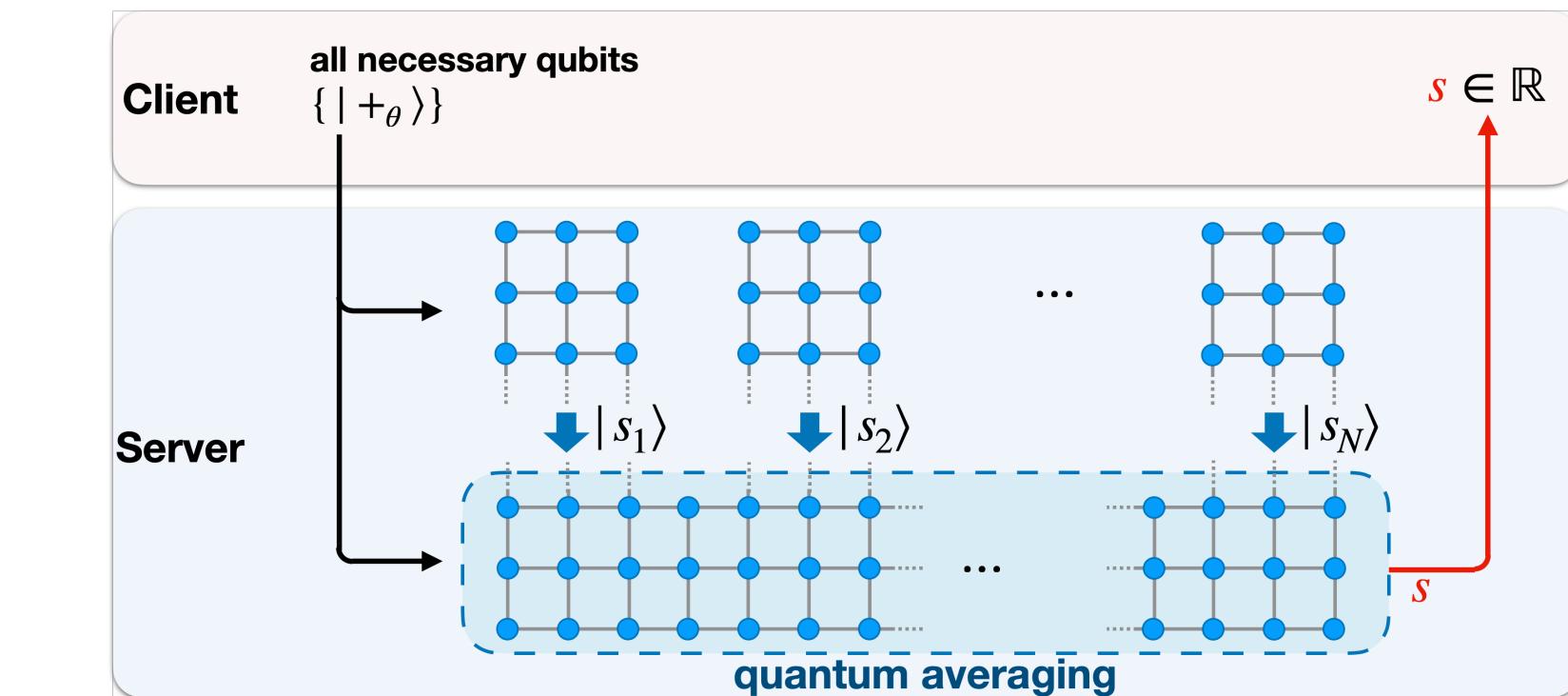
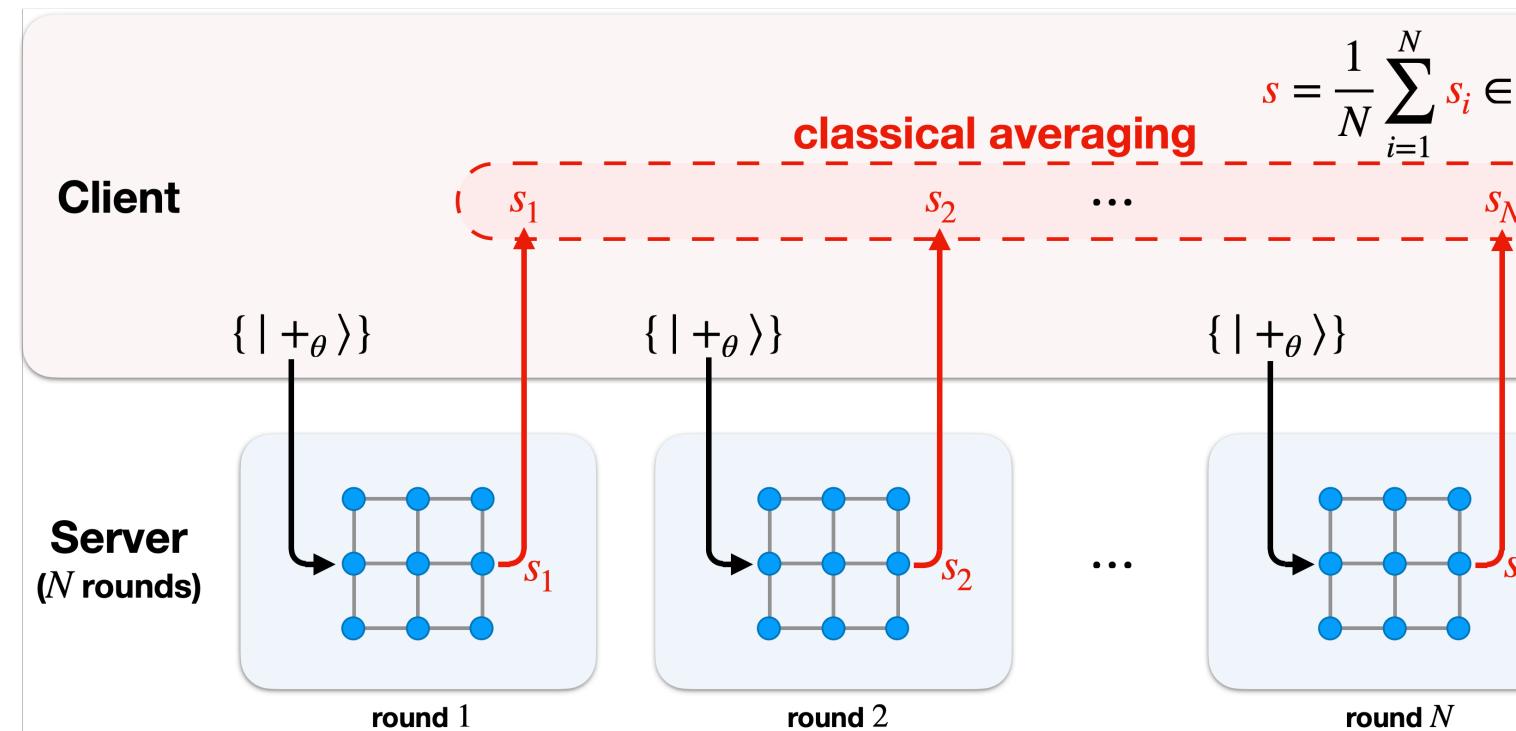
↓ Difference in Accept/Abort probability

Distinguishability: $O (e^{-N_c} + e^{-N_t})$

Impact of this work

Bridge between crypto-protocols and **key quantum utility**

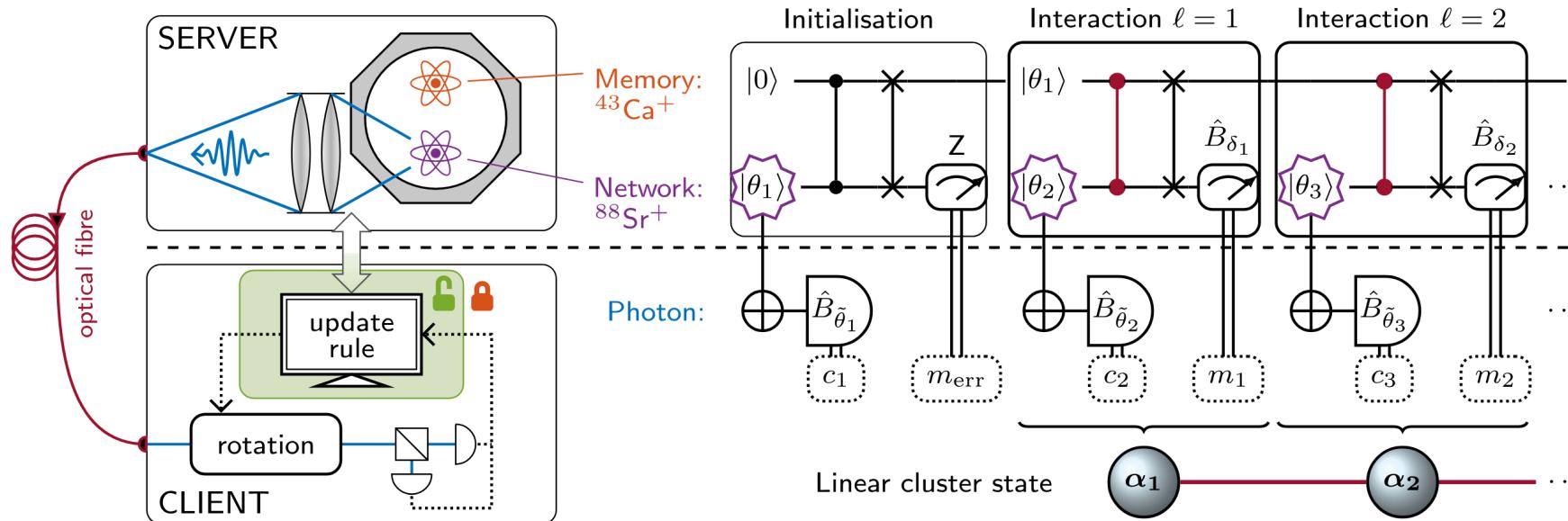
- A new avenue for cryptographic **security analysis** of near-term quantum-classical hybrid **techniques with expectation values**
- The only known efficient **overhead-free** verification of near-term tasks with potential **quantum advantage**
- A convergence between foundational **cryptographic theory** and **physically motivated quantum tasks**



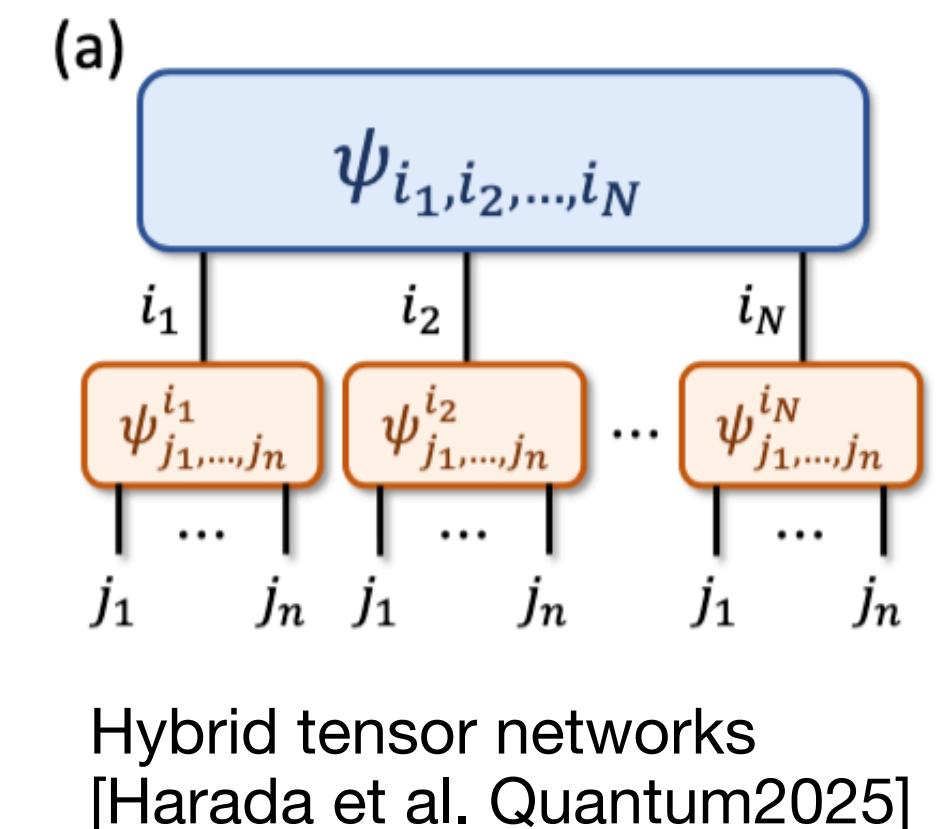
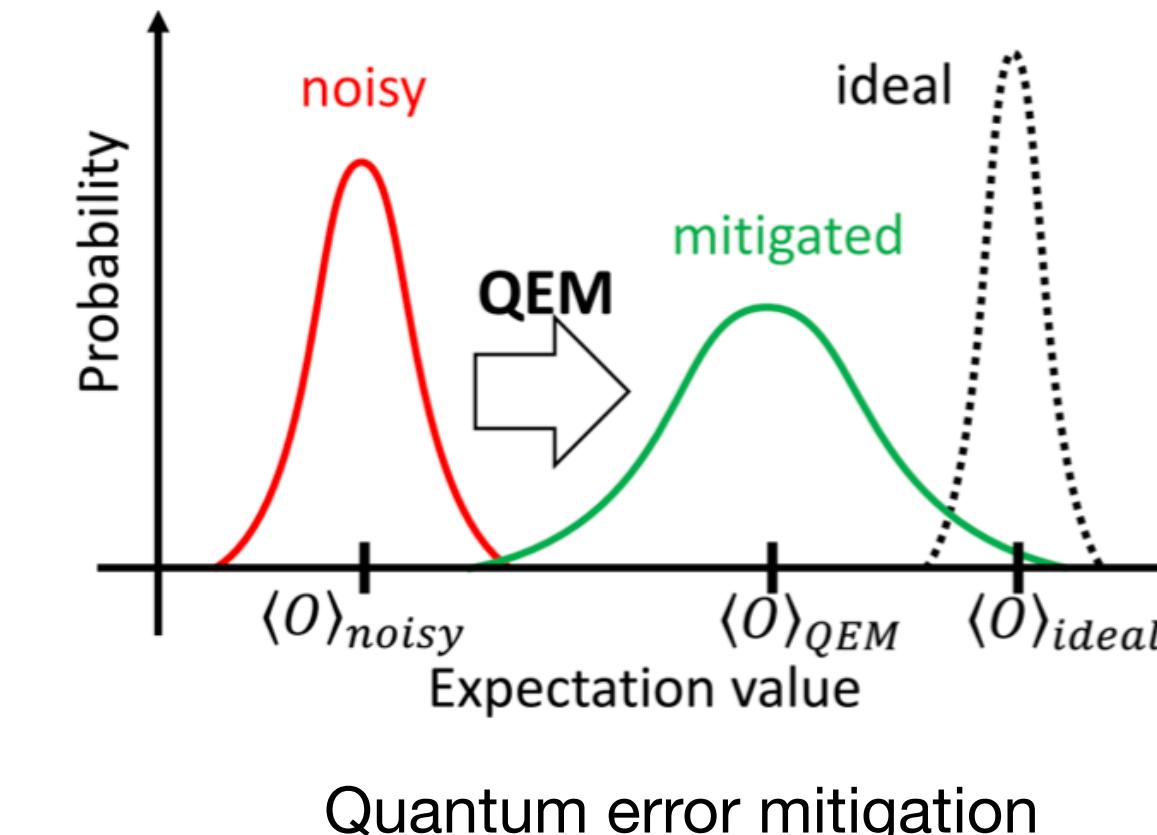
Future work

A broad class of applications of promising interest

- Integration with **quantum error mitigation** and quantum-classical hybrid approaches
- Verifiable quantum advantage **experiments** on current hardware



Implementation in trapped-ion devices [Drmota et al. PRL2024]

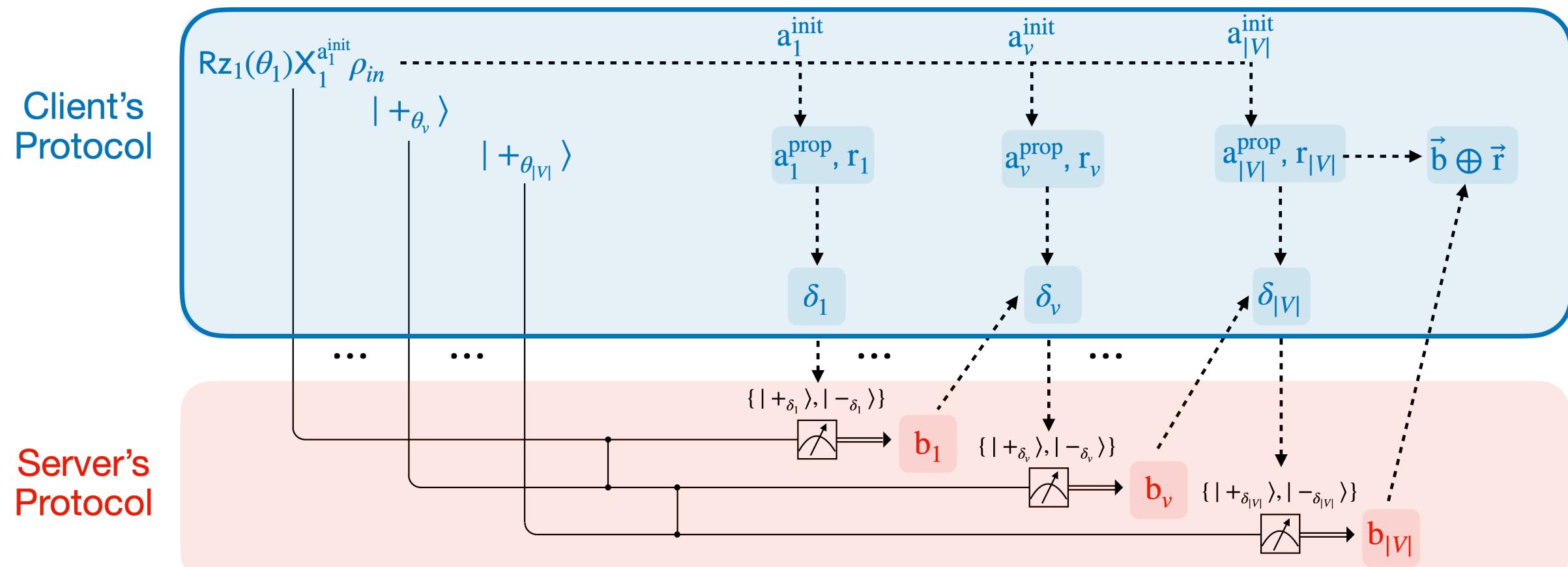


Link to our paper: <https://www.arxiv.org/abs/2510.08548>

Appendix

<https://scirate.com/arxiv/2510.08548>

Separating Client's protocol of VBOE into Resource Part and Simulator Part



VBOE protocol

